Networked Distributed Energy Resources

Steven Low

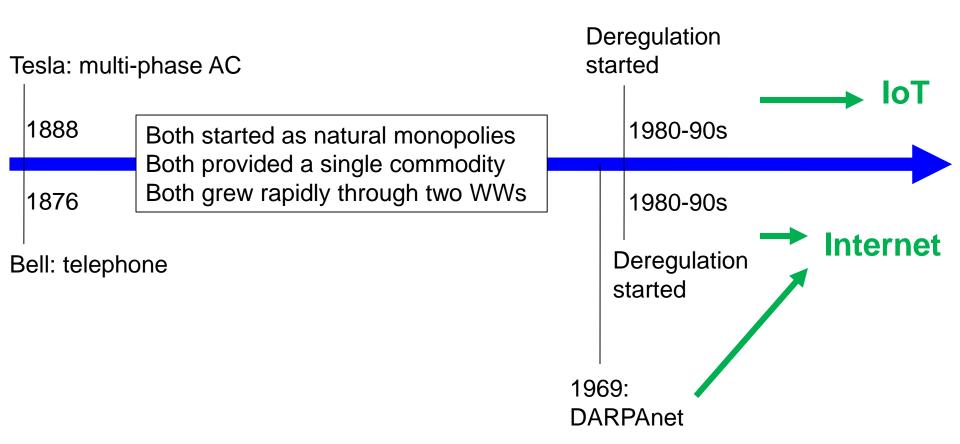
Electrical Engineering Computing + Math Sciences



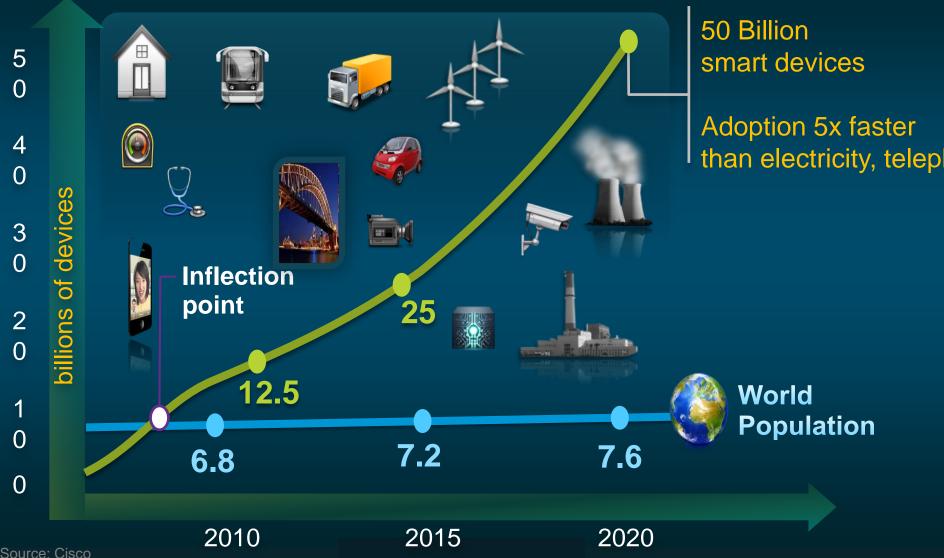
Sept 2015 Energy Colloquium, Skoltech, Moscow



Power network will undergo similar <u>architectural</u> <u>transformation</u> that phone network went through in the last two decades



Internet of Things (IoT)



Source: Cisco IBSG, 2011



Industries will be destroyed and created AT&T, MCI, McCaw Cellular, Qualcom Google, Facebook, Twitter, Amazon, eBay, Netflix

Infrastructure will be reshaped

Centralized intelligence, vertically optimized Distributed intelligence, layered architecture

What will drive power network transformation ?



Proliferation of renewables

Electrification of transportation

Advances in power electronics

Deployment of sensing, control, comm

challenge

enabler



Big picture Energy stats and trends

Challenges and opportunities Implications on smart grid

Sample Caltech research

- Control and optimization
- Power flow and dynamics





World energy stats (2011)

Consumption	519 quad BTU
petroleum	34%
coal	29%
gas	23%
renewable (elec)	8%
nuclear	5%

	Consumption	519 (quad BTU)	per capita (mil BTU)
top 5 countries	China	20%	78
	US	19%	313
	Russia	6%	209
	India	5%	20
	Japan	4%	164
	total	54%	0

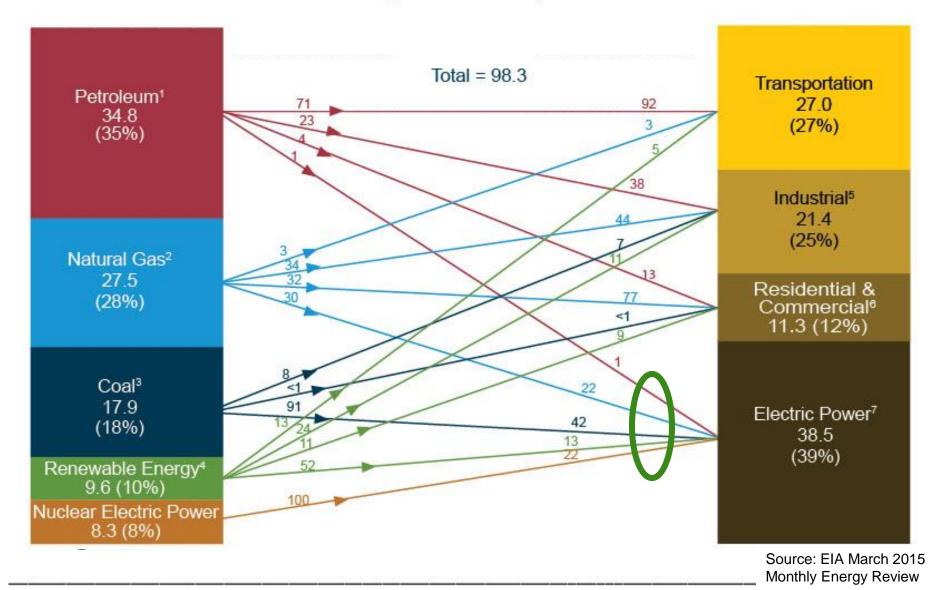
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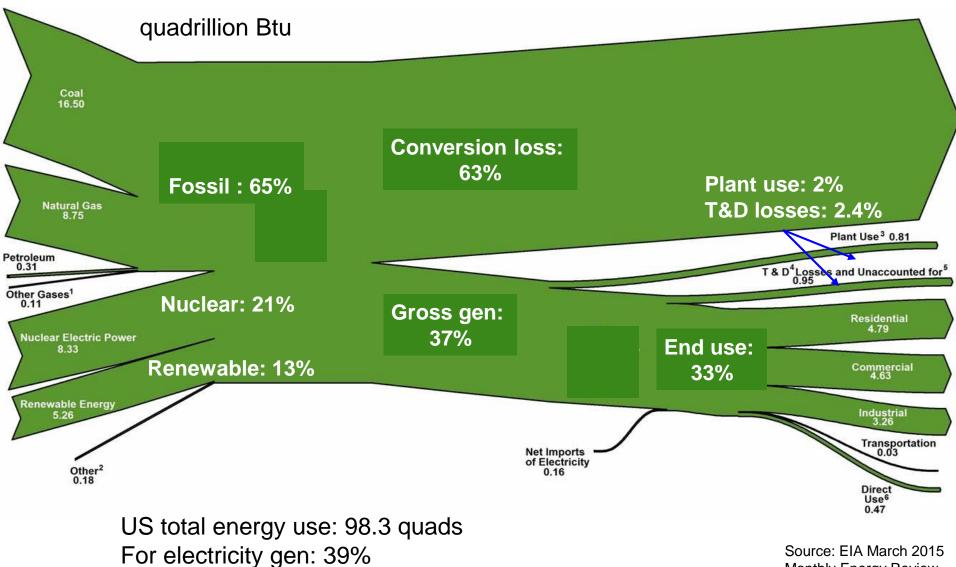
	Consumption	519 (quad BTU)	CO2 emission
top 5 countries	China	20%	27%
	US	19%	17%
	Russia	6%	5%
	India	5%	5%
	Japan	4%	4%
	total	54%	58%

US Primary Energy Flow 2014

(Quadrillion Btu)

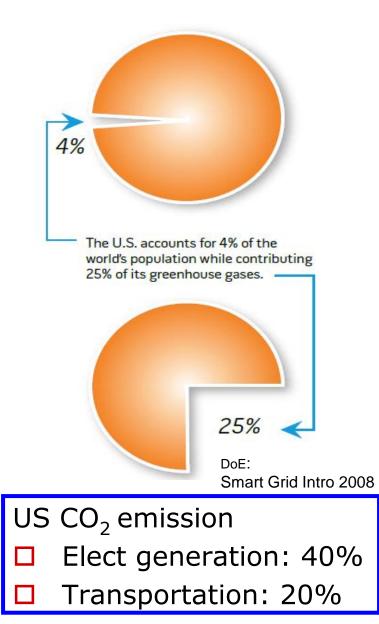


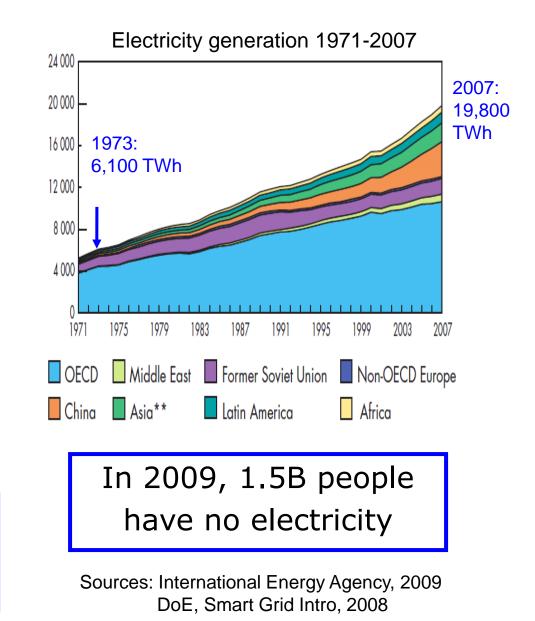




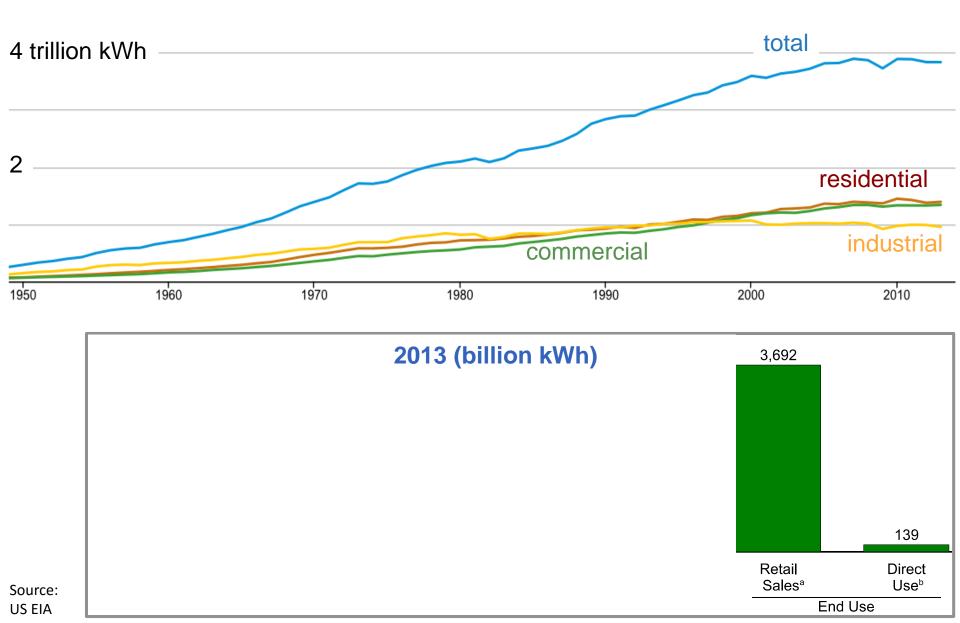
Monthly Energy Review

Sustainability challenge

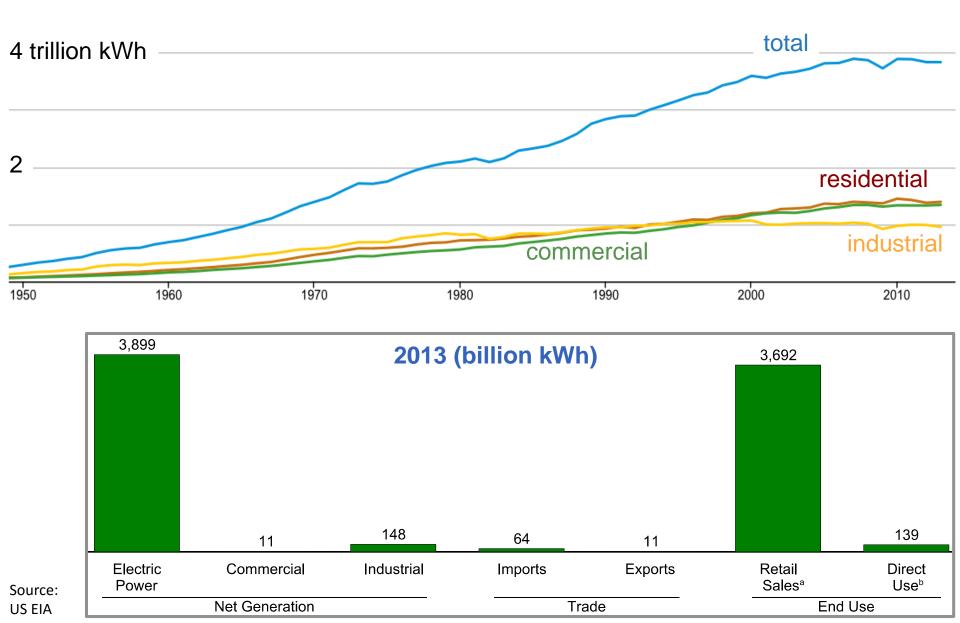


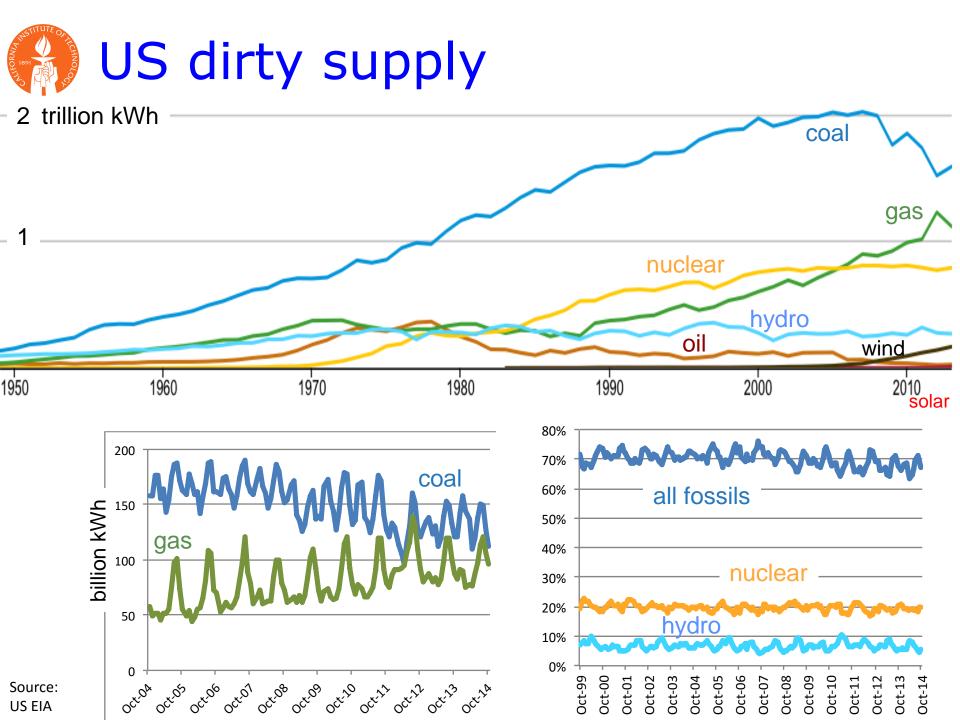




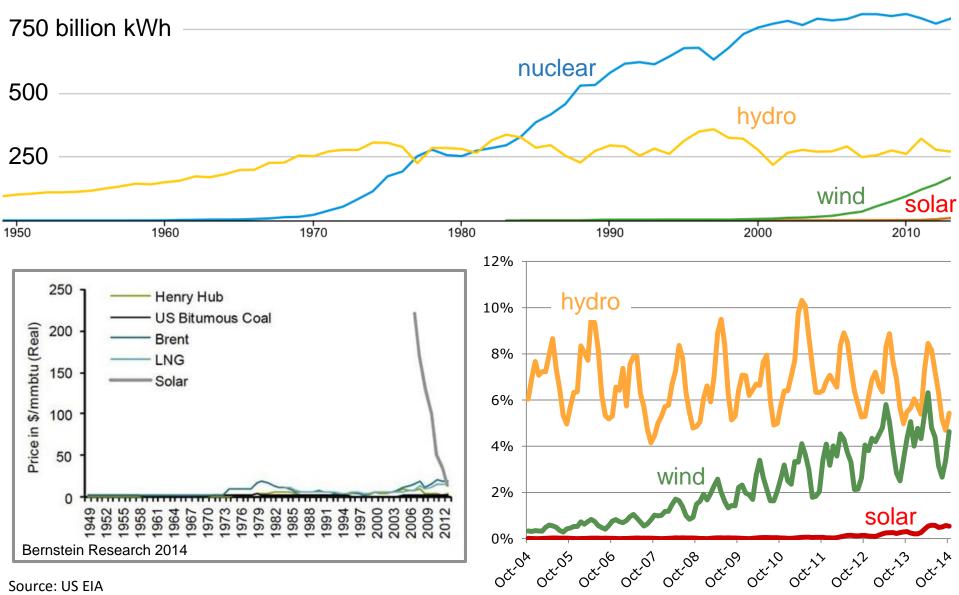




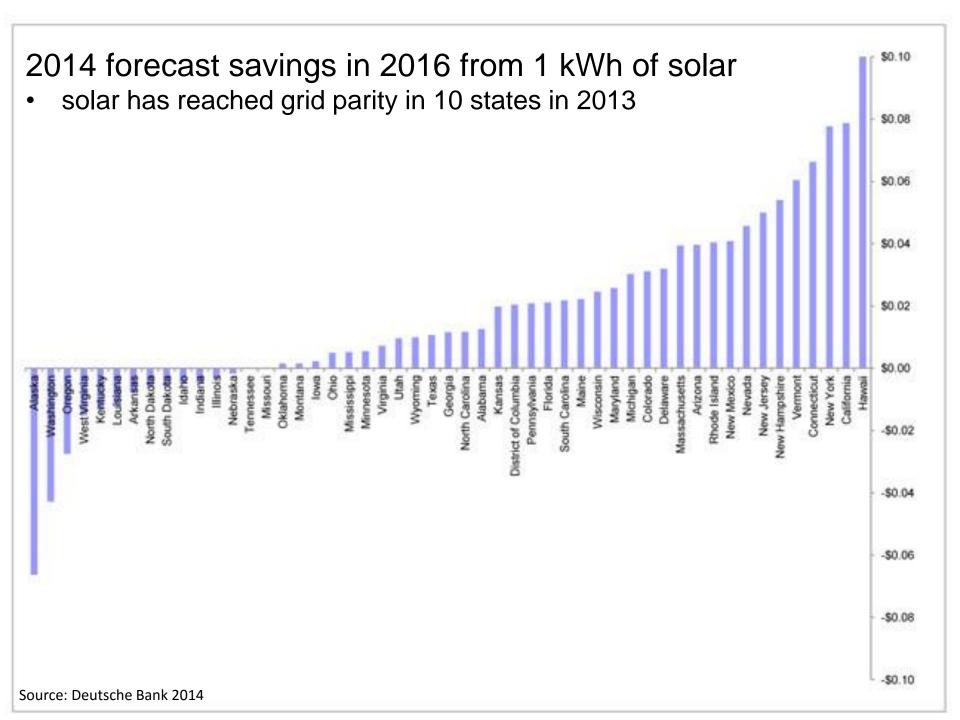




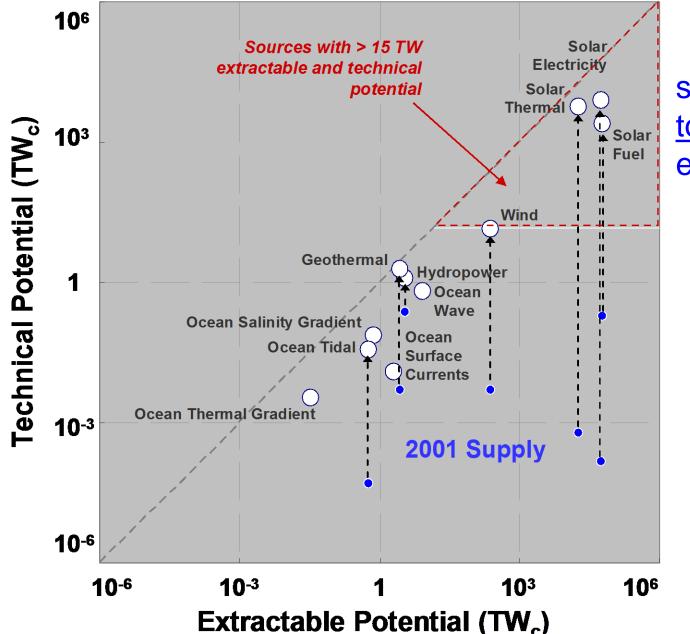
US renewable generations CHNOZ,



Source: US EIA



Technical potential of renewable sources



solar can supply total world energy demand

> Source: Tsao, Lewis, Crattree, 2006

Area to power the world by solar

1980 (based on actual use) 207,368 SQUARE KILOMETERS

2008 (based on actual use) 366,375 SQUARE KILOMETERS

2030 (projection) 496,805 SQUARE KILOMETERS

- Areas are calculated based on an assumption of 20% operating efficiency of collection devices and a 2000 hour per year natural solar input of 1000 watts per square meter striking the surface.
 - These 19 areas distributed on the map show roughly what would be a reasonable responsibility for various parts of the world based on 2009 usage. They would be further divided many times, the more the better to reach a diversified infrastructure that localizes use as much as possible.
 - The large square in the Saharan Desert (1/4 of the overall 2030 required area) would power all of Europe and North Africa. Though very large, it is 18 times less than the total area of that desert.
 - The definition of "power" covers the fuel required to run all electrical consumption, all machinery, and all forms of transportation. It is based on the US Department of Energy statistics of worldwide Btu consumption and estimates the 2030 usage (678 quadrillion Btu) to be 44% greater than that of 2008.
 - Area calculations do not include magenta border lines.



Big pictureEnergy stats and trends

Challenges and opportunities Implications on smart grid

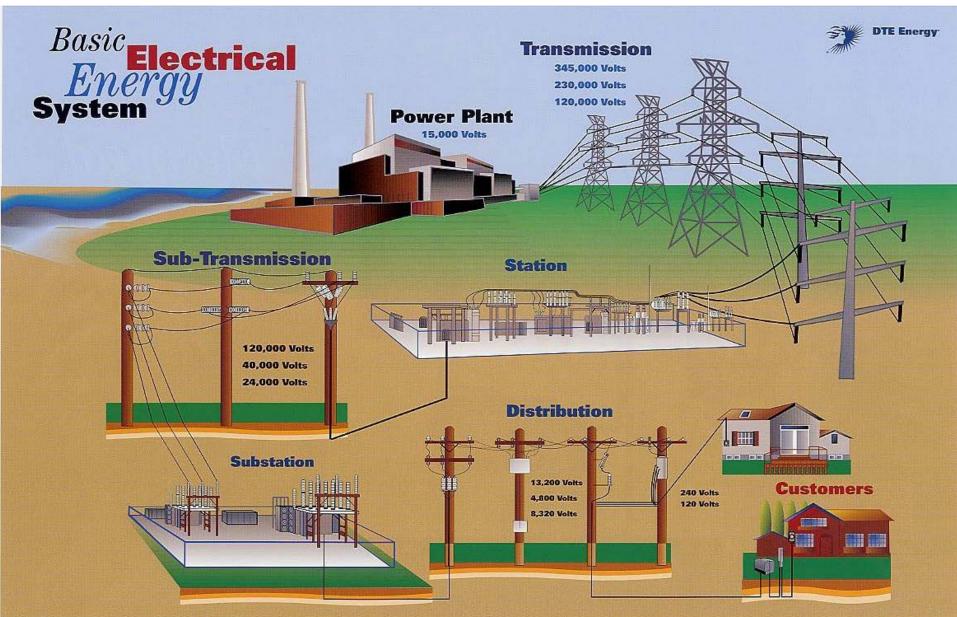
Sample Caltech research

- Control and optimization
- Power flow and dynamics

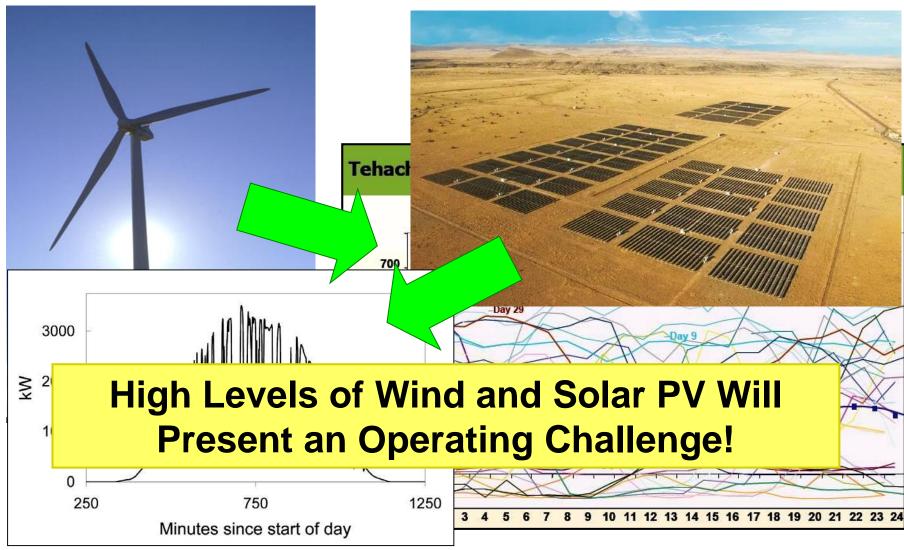






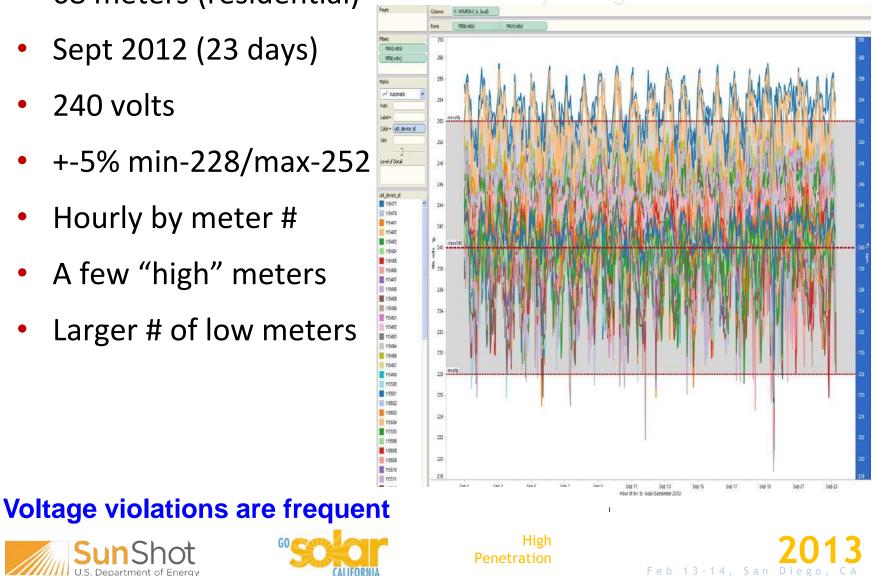






Source: Rosa Yang, EPRI

- 68 meters (residential)
- Sept 2012 (23 days)
- 240 volts
- +-5% min-228/max-252
- Hourly by meter #
- A few "high" meters
- Larger # of low meters



Source: Leon Roose, University of Hawaii Development & demo of smart grid inverters for high-penetration PV applications



with the issues fast.

"We have here a graving morther of contention building for all gold solutions," sold (Sels Deliver, managing parties of Karnel R and Harvel Range Commissio.

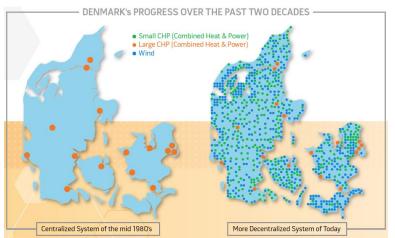


1 Proliferation of renewables

- Driven by sustainability
- Enabled by policy and investment

2 Migration to distributed arch

- 2-3x generation efficiency
- Relief demand on grid capacity

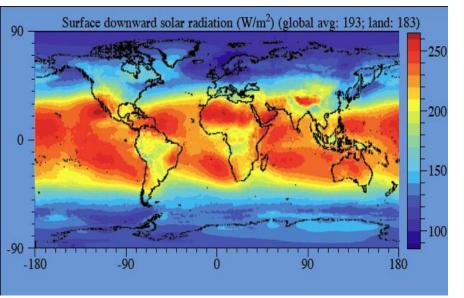




1 Proliferation of renewables

- Driven by sustainability
- Enabled by policy and investment
- 2 Migration to distributed arch
 - 2-3x generation efficiency
 - Relief demand on grid capacity
- 3 Rise of Internet of Things (IoT)
 - 5x adoption rate of electricity & phone
 - Impact on industry and residential

Solar power over land: > 20x world energy demand







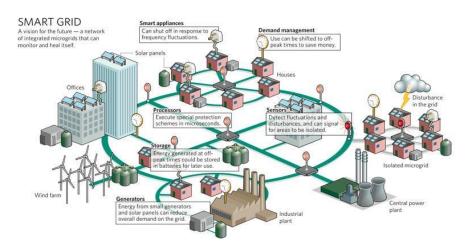
network of billions of active distributed energy resources (DERs)

DER: PV, wind tb, EV, storage, smart bldg / appl

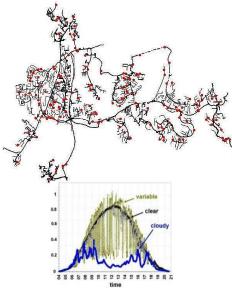
Risk: active DERs introduce rapid random fluctuations in supply, demand, power quality increasing risk of blackouts



Opportunity: active DERs enables realtime dynamic network-wide feedback control, improving robustness, security, efficiency



Caltech research: distributed control of networked DERs



- Foundational theory, practical algorithms, concrete applications
- Integrate engineering and economics
- Active collaboration with industry













Current control paradigm works well today

- Centralized, open-loop, human-in-loop, worst-case preventive
- Low uncertainty, few active assets to control
- Schedule supplies to match loads

Future needs

- Closing the loop, e.g. real-time DR, Volt/VAR control, EV/storage mgt
- Fast computation to cope with rapid, random, large fluctuations in supply, demand, voltage, freq
- Simple algorithms to scale to large networks of active DERs



intelligence everywhere connected



Global energy demand will continue to grow

Traditional supply is unsustainable

There is more renewable energy than the world ever needs

Someone will figure out how to capture and store it

There will be connected intelligence everywhere

- Cost of computing, storage, communication and manufacturing will continue to drop
- Power system will transform into the largest and most complex Internet of Things
 - Generation, transmission, distribution, consumption, storage



To develop technologies that will enable and guide the historic transformation of our power system

- Generation, transmission, distribution, consumption, storage
- Devices, systems, theory, algorithms
- Control, optimization, stochastics, data, economics



Large scale

Distributed algorithms

Uncertainty

Risk-limiting approach

Multiple timescales

Decomposition

Nonconvexity

Convex relaxations



Example: Southern California Edison

4-5 million customers

SCE Rossi feeder circuit

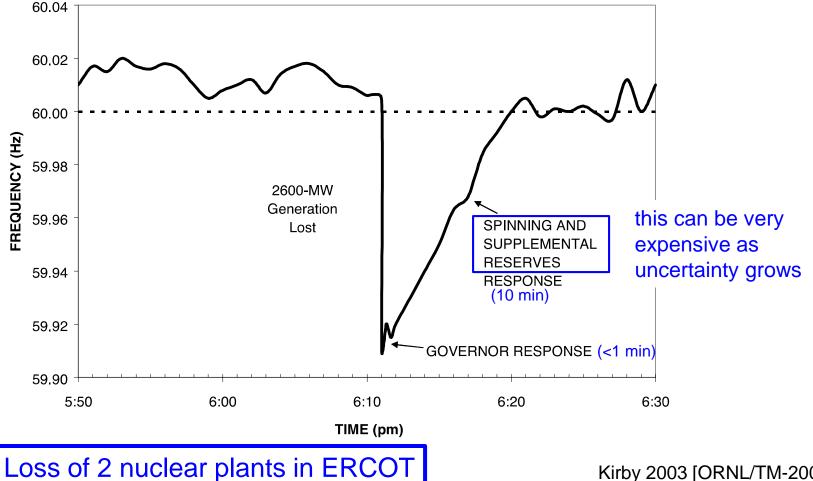
- #houses: 1,407; #commercial/industrial: 131
- #transformers: 422
- #lines: 2,064 (multiphase, inc. transformers)
- peak load: 3 6 MW
- #optimization variables: 50,000
- SCE has 4,500 feeders
 - ~100M variables

United States

 131M customers, 300K miles of transmission & distr lines, 3,100 utilities



Uncertainty creates difficulty in both control and markets



Kirby 2003 [ORNL/TM-2003/19]



Real-time price can be more than 100x the average price !

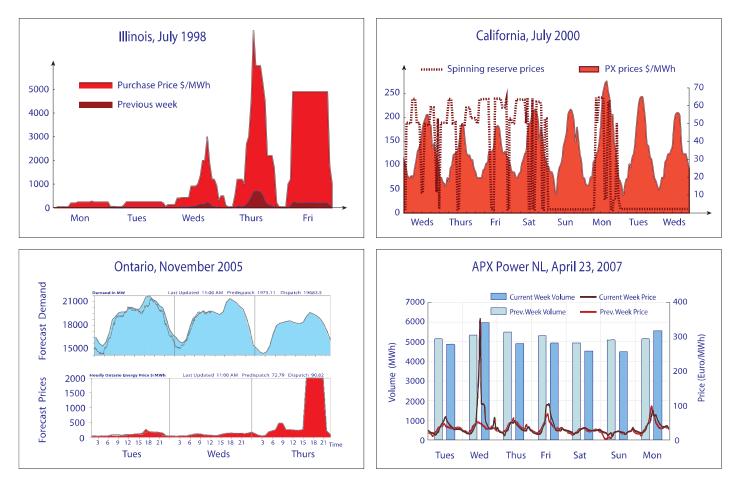


Figure: Real-world price dynamics < a > < = > <

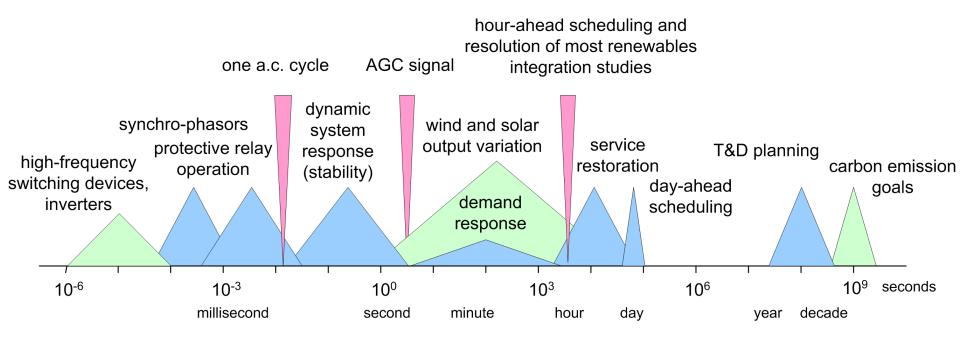
Sean Meyn, 2010

Ξ



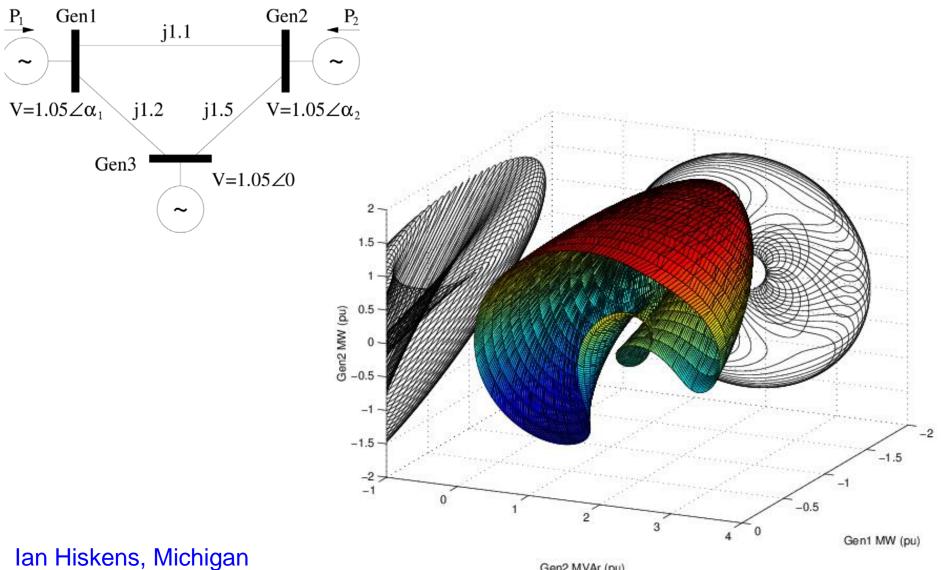
System dynamics and controls at different timescales

- require different models
- they interact



Sean Meyn, 2010





Gen2 MVAr (pu)



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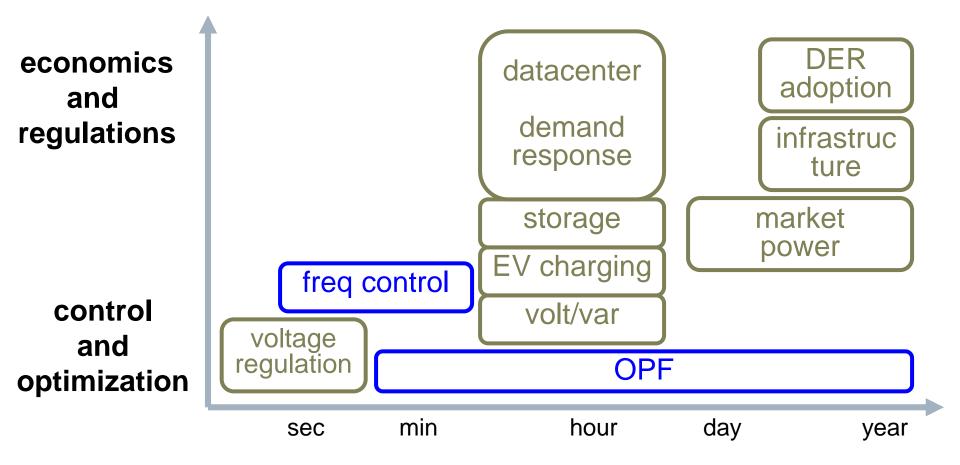
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6+ faculty

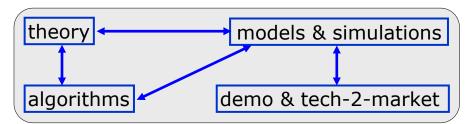


optimal power flow



Distributed Control of Networked DER an Groad GENI project





Caltech: Profs Chandy, Doyle, Low (PI); Drs. Bunn, Mallada; Students: Agarwal, Cai, Chen, Farivar, Gan, Guo, Matni, Peng, Ren, Tang, You, Zhao SCE: Auld, Castaneda, Clarke, Gooding, Montoya, Shah, Sherick (PI) Newport/Caltech: DeMartini (advisor)

Alumni: Bose (Cornell), Chen (Colorado), Collins (USC), Gayme (JHU), Lavaei (Columbia), Li (Harvard), Topcu (UPenn), Xu (SUTD)

quadratic in V

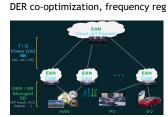
linear in W !

nonconvex constr

EAN

- Increase(asset(u+liza+on(and(efficiency(
- Improve(power(quality(and(stability(
- Move(data:in:mo+on(analy+cs(to(edge(

Contact: Michael Enescu, co-founder CEO, enescu@alumni.caltech.edu



• EAN analytics and optimization

EAN enabled control

DER placement, asset opt, analytics

applications and T2M

theory

Convex relaxation of OPF:

OPF: min tr (CVV^*)

min tr (CW)

SDP relaxation

Theoretical foundation for semi-

definite relaxations of power flow

s. t. $\underline{s}_i \notin \operatorname{tr}(Y_i^* V V_{\bullet}^*) \notin \overline{s}_i, \ \underline{v}_i \notin |V_i|^2 \notin \overline{v}_i$

s.t. $s_i \in \operatorname{tr}(Y_i^*W) \in \overline{s}_i, \quad V_i \in W_{ii} \in \overline{V}_i$

Exact relaxations: Sufficient conditions for recovering global

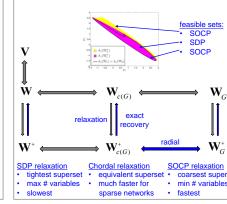
optimum of OPF from relaxations

 W^{3} 0, rank W = 1 ignore this (only)

algorithms

Relaxation algorithms:

- single-phase balanced, multiphase unbalanced
- centralized, distributed



models

DER adoption model & software Sophisticated feedback model • Cloud service for PV-uptake: http://etechuptake.appspot.com/ EDISON'

Increased T&D infrastructure

volt/var control with renewables

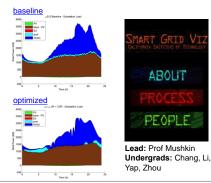
SCE circuits, DER forecasts

advanced OPF solver

simulations

Realistic simulations

- SCE feeder model, 2,000 buses
- DER: inverters, HVAC, pool pumps, EV
- Multiphase unbalanced radial





OPF is solved routinely for

- network control & optimization decisions
- market operations & pricing
- at timescales of mins, hours, days, ...

Non-convex and hard to solve

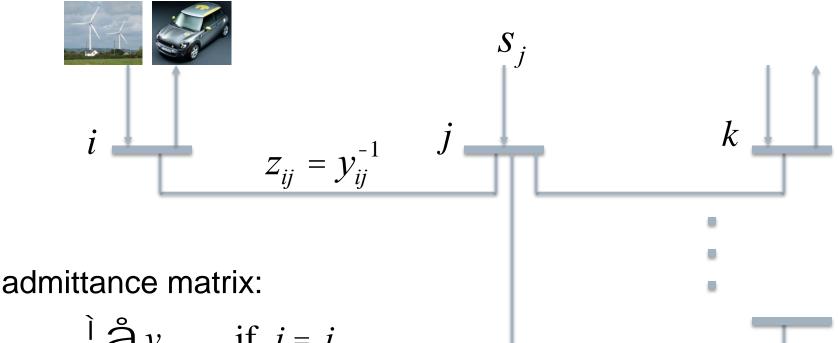
- Huge literature since 1962
- Common practice: DC power flow (LP)
- Also: Newton-Raphson, interior point, ...



OPF underlies many applications

- Unit commitment, economic dispatch
- State estimation
- Contingency analysis
- Feeder reconfiguration, topology control
- Placement and sizing of capacitors, storage
- Volt/var control in distribution systems
- Demand response, load control
- Electric vehicle charging
- Market power analysis





 $Y_{ij} := \begin{bmatrix} i & a & y_{ik} & \text{if } i = j \\ i & -y_{ij} & \text{if } i \sim j \\ i & 0 & \text{else} \end{bmatrix}$

graph G: undirected

 \boldsymbol{Y} specifies topology of \boldsymbol{G} and impedances \boldsymbol{z} on lines



In terms of V:

$$S_j = \operatorname{tr}\left(Y_j^H V V^H\right)$$
 for all j $Y_j = Y^H e_j e_j^T$

Power flow problem: Given (Y, s) find V



isolated solutions



mintr (CVV^H) gen cost,
power lossover(V, s)subject to \underline{s}_j for s_j for \overline{s}_j \underline{V}_j for V_j for V_j for V_j



mintr
$$(CVV^H)$$
gen cost,
power lossover (V,s) subject to \underline{s}_j f s_j f \overline{s}_j \underline{V}_j f $|V_j|$ f \overline{V}_j $s_j = \operatorname{tr} (Y_j^H VV^H)$ power flow equation



min
$$\operatorname{tr} CVV^H$$

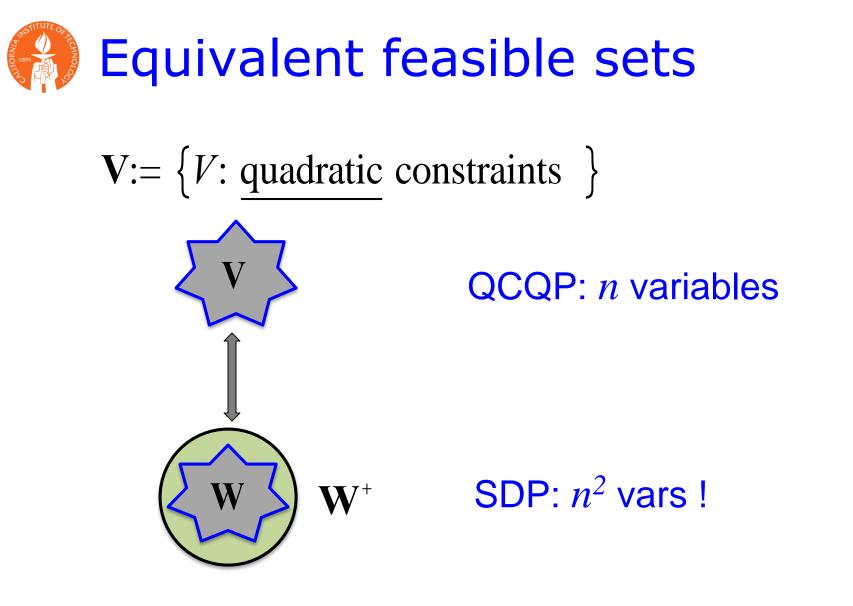
subject to \underline{s}_j \pounds $\operatorname{tr} \left(Y_j VV^H \right)$ \pounds \overline{s}_j \underline{v}_j \pounds $|V_j|^2$ \pounds \overline{v}_j

nonconvex QCQP (quad constrained quad program)



min tr
$$CVV^H$$

subject to $\underline{s}_j \in \operatorname{tr}(Y_jVV^H) \in \overline{s}_j \quad \underline{v}_j \in |V_j|^2 \in \overline{v}_j$
quadratic in V
linear in W
subject to $\underline{s}_j \in \operatorname{tr}(Y_jW) \in \overline{s}_j \quad \underline{v}_i \in W_{ii} \in \overline{v}_i$
 $W^3 0, \operatorname{rank} W = 1$ convex in W
except this constraint



W:= {W: linear constraints } $\bigcap \{W \ge 0 \text{ rank-1}\}$ idea: $W = VV^H$



$$\mathbf{W}_G := \left\{ W_G : \underline{\text{linear}} \text{ constraints} \right\}$$

idea:
$$W_G = (VV^H \text{ only on } G)$$

W:= {W: linear constraints }
$$\bigcap \{W \ge 0 \text{ rank-1}\}$$

idea: $W = VV^H$



$$\mathbf{W}_{G} := \left\{ W_{G} : \underline{\text{linear}} \text{ constraints } \right\} \cap \left\{ \begin{aligned} W(j,k) \ge 0 \text{ rank-1,} \\ \text{cycle cond on } \angle W_{jk} \end{aligned} \right\}$$

idea: $W_{G} = \left(VV^{H} \text{ only on } G \right)$

$$W:= \{W: \underline{\text{linear} \text{ constraints}}\} \cap \{W \ge 0 \text{ rank-1}\}$$

idea: $W = VV^H$





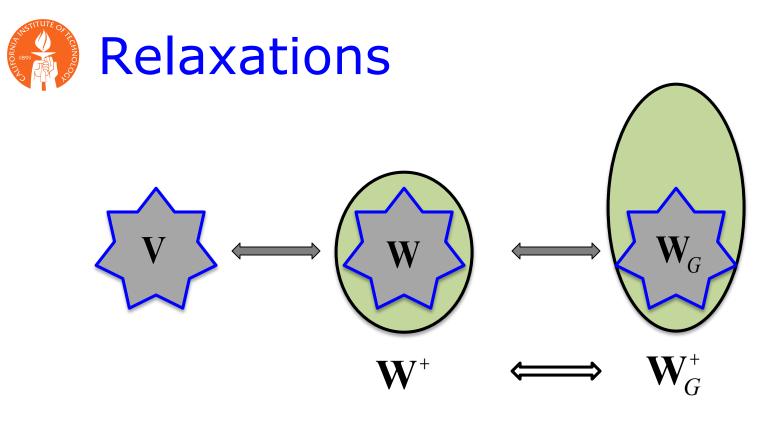
Theorem: $\mathbf{V}^{\circ}\mathbf{W}^{\circ}\mathbf{W}_{G}$

Bose, Low, Chandy Allerton 2012 Bose, Low, Teeraratkul, Hassibi TAC2014



$$\mathbf{W}_{G} := \left\{ W_{G} : \underline{\text{linear}} \text{ constraints } \right\} \cap \left\{ \begin{matrix} W(j,k) \ge 0 \text{ rank-1}, \\ \text{cycle cond on } \angle W_{jk} \end{matrix} \right\}$$

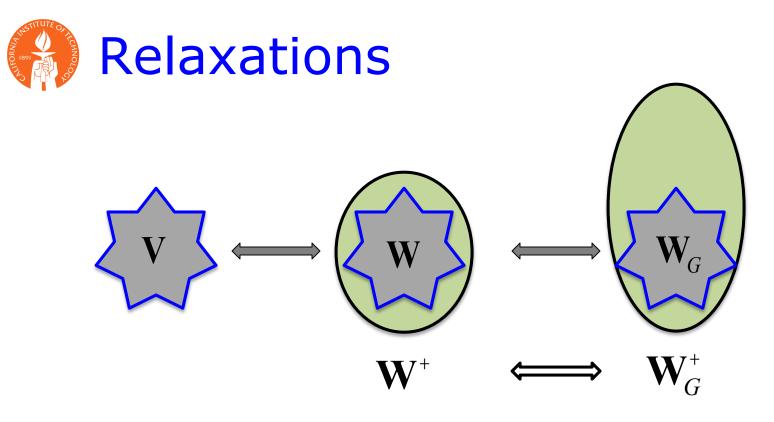
$$\mathbf{W} := \left\{ W : \underline{\text{linear}} \text{ constraints } \right\} \cap \left\{ W \ge 0 \text{ rank-1} \right\}$$



Theorem

- **Radial** G : $\mathbf{V} \subseteq \mathbf{W}^+ @ \mathbf{W}_G^+$
- $\blacksquare \mathsf{Mesh} G : \mathsf{V} \subseteq \mathsf{W}^{\scriptscriptstyle +} \subseteq \mathsf{W}_{G}^{\scriptscriptstyle +}$

Bose, Low, Chandy Allerton 2012 Bose, Low, Teeraratkul, Hassibi TAC2014



<u>Theorem</u>

- **Radial** G : $\mathbf{V} \subseteq \mathbf{W}^+ @ \mathbf{W}_G^+$
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For radial networks: always solve SOCP !



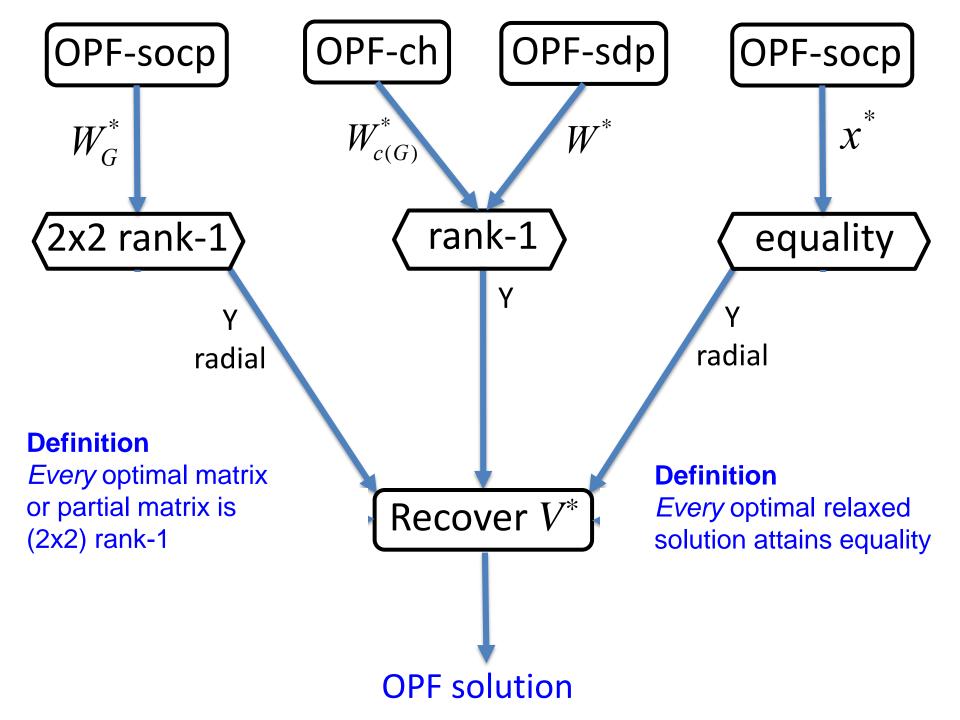
OPF $\min_{V} C(V) \text{ subject to } V \mid \mathbf{V}$

OPF-sdp:

$\min_{W} C(W_G) \quad \text{subject to} \quad W \in \mathbb{W}^+$

OPF-socp:

$\min_{W_G} C(W_G) \quad \text{subject to} \quad W_G \in W_G^+$



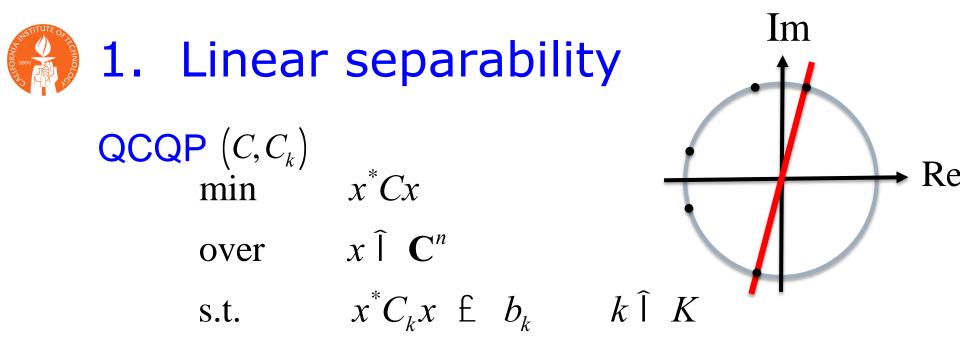


$\begin{array}{ccc} \mathsf{QCQP} (C, C_k) \\ \min & x^* C x \\ \text{over} & x \widehat{|} \mathbf{C}^n \\ \text{s.t.} & x^* C_k x \widehat{|} b_k & k \widehat{|} K \end{array}$

graph of QCQP

 $G(C, C_k)$ has edge (i, j) $C_{ij} \Box 0$ or $[C_k]_{ij} \Box 0$ for some k

QCQP over tree $G(C, C_k)$ is a tree

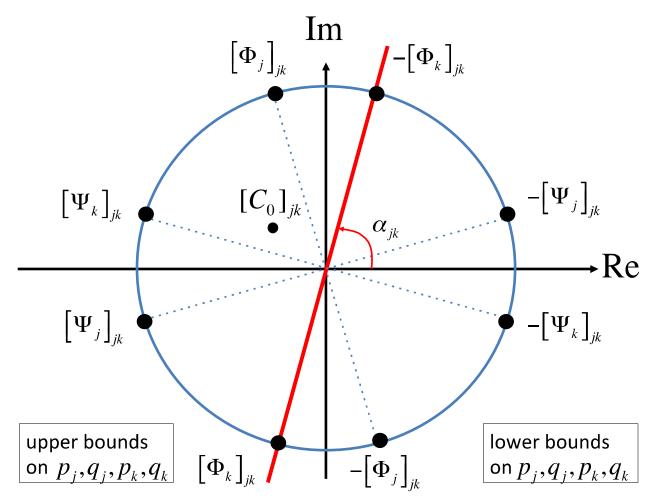


Key condition
$$i \sim j: (C_{ij}, [C_k]_{ij}, "k)$$
 lie on half-plane through 0

Theorem SOCP relaxation is exact for QCQP over tree

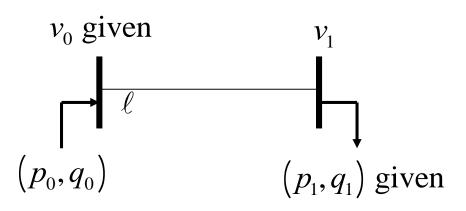
Bose et al 2012 Sojoudi, Lavaei 2013

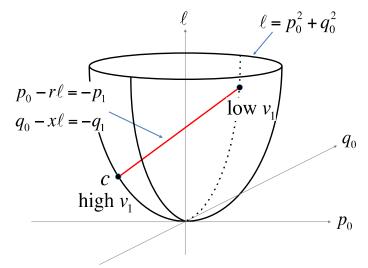




Not both lower & upper bounds on real & reactive powers at both ends of a line can be finite

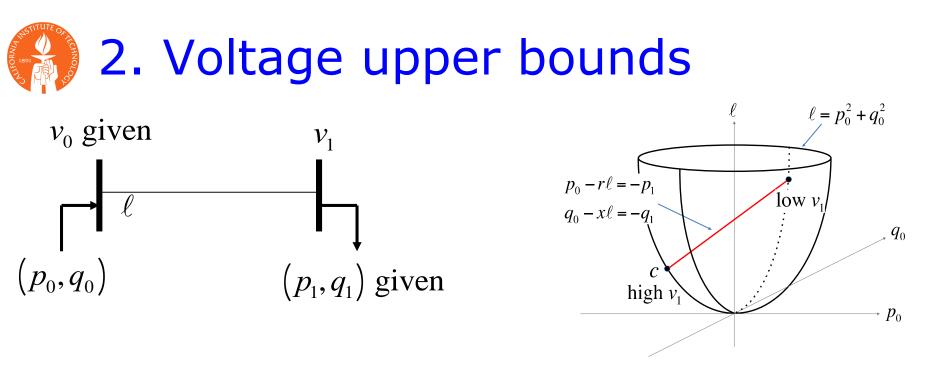




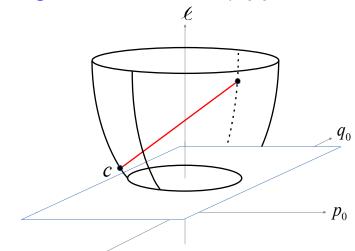


when there is no voltage constraint

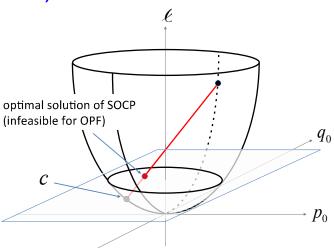
- feasible set : 2 intersection pts
- relaxation: line segment
- exact relaxation: c is optimal



voltage lower bound (upper bound on l) does not affect relaxation



(a) Voltage constraint not binding



(b) Voltage constraint binding



OPF:
$$\min_{x \in \mathbf{X}} f(x)$$
 s.t. $\underline{v} \in v \in \overline{v}, s \in \mathbf{S}$

SOCP: $\min_{x \in \mathbf{X}^+} f(x)$ s.t. $\underline{v} \in v \in \overline{v}$, $s \in \mathbf{S}$

Key condition:

- $L(s) \notin \overline{v}$
- Jacobian condition $\underline{A}_{i_t} \cdots \underline{A}_{i_t 0} Z_{i_t 0, 1} > 0$ for all $1 \le t \le t^0 < k$

voltages if network were lossless

if upward current were reduced then all subsequent powers dec

Theorem SOCP relaxation is exact for radial networks

Gan, Li, Topcu, Low TAC2014



OPF:
$$\min_{x \in \mathbf{X}} f(x)$$
 s.t. $\underline{v} \in v \in \overline{v}, s \in \mathbf{S}$

SOCP: $\min_{x \in \mathbf{X}^+} f(x)$ s.t. $\underline{v} \in v \in \overline{v}$, $s \in \mathbf{S}$

Key condition:

- $L(s) \notin \overline{v}$
- Jacobian condition $\underline{A}_{i_t} \cdots \underline{A}_{i_t 0} Z_{i_t 0+1} > 0$ for all $1 \le t \le t^0 < k$

satisfied with large margin in IEEE circuits and SCE circuits

Theorem SOCP relaxation is exact for radial networks

Gan, Li, Topcu, Low TAC2014



